

**PLEASE RETURN THIS SHEET ALONG WITH ALL
THE SHEETS YOU WERE GIVEN**

SURNAME _____

1) (POINTS 8/30) Consider a bus-based multicore that supports a new cache-coherence protocol called MESIF. Compared to the well-known MESI protocol, the MESIF protocol adds a 5th state called F (Forward). A copy in F state is like an S copy but the cache which holds it has the responsibility to provide (forward) the copy once a BusRd transaction involves that copy; at the same time, the S state is now simplified, as it doesn't have to respond to cache-to-cache transfers (Flush* transactions) due to a BusRd. There will be at most one cache holding a copy in the F state; on a PrRd miss, if there are other shared copies, the state will be F, whilst if another F copy observes the BusRd will go to S. Both F and S states are clean shared states. If the copy in F state is eventually evicted, there will be no copy in the F state: in this case the copy will be provided by the memory.

1a [Points 8/30] Draw the diagram of the MESIF protocol according to the above description.

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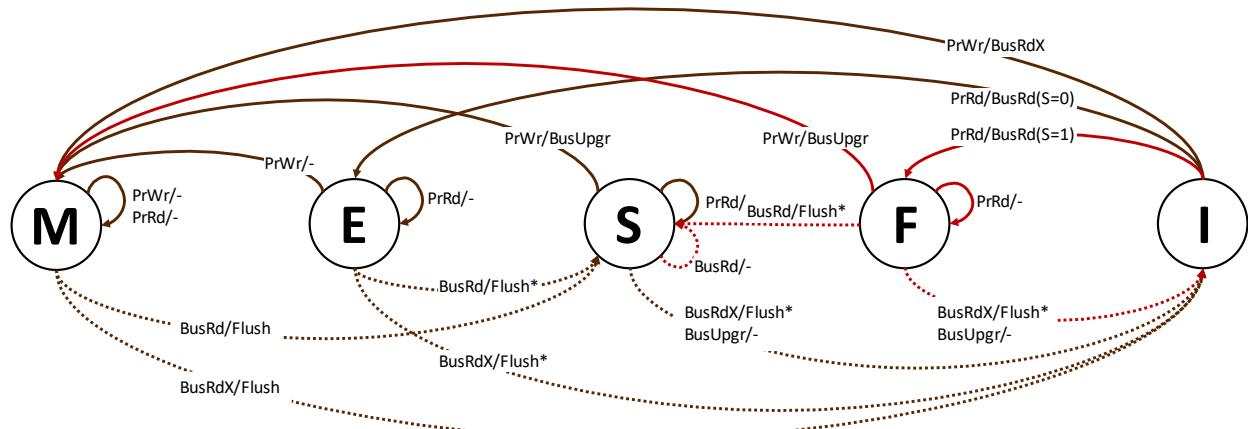
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1b) [Points 22/30] Assuming a cost of 1cc (1 clock-cycle) for read/write operations, 90cc for BusRd or BusRdx transactions, 60cc for BusUpgr, 20 cc for Flush* and 30cc for Flush. Evaluate the total cost (in clock-cycles) for the following streams:

EXERCIZE 1a)

First, we start drawing the states of the MESI protocol (except for S state – see brown lines), then let's focus on S and F states. **S-state:** PrRd and PrWr are exactly as in MESI; also, BusRdX and BusUpgr, as in MESI, cause an invalidation of the local copy; BusRdX will also provide a copy (useful, e.g., in the case of write-miss (PrWr from I)) to be written partially; differently from MESI, now a BusRd - snooped from a cache with a copy in S-state – will NOT be followed by a Flush*: the cache with the shared copy in F-state will provide such copy, eventually – if such copy has been evicted in the meantime, the memory will intervene.

F: state: PrRd and PrWr are exactly like for the S-state (remember that F is a shared clean state); like in S-state, the BusRdX and BusUpgr transaction cause a local invalidation, with the Flush* following the BusRdX; we enter the F-state on a read-miss once the shared-line indicates that there exist other copies in the system (S=1), so the newly acquired copy is shared; when a BusRd is snooped, this cache has the responsibility to provide the copy via a Flush* and its state will be then demoted to S (its duty is finished, the other cache with the shared copy will get the new copy in F state – it is the only possibility for a BusRd – the other is with S=0, which is not this case).



EXERCIZE 1b)

Stream-1 MESIF	Core Operation	C1	C2	C3	Bus Transaction	Data from	Cycles
	PrRd1	E	I	I	BusRd(S=0)	Mem	90
	PrWr1	M			-	-	1
	PrRd1	M			-	-	1
	PrWr1	M			-	-	1
	PrRd2	S	F		BusRd(S=1), Flush	C1	90+30
	PrWr2	I	M		BusUpgr	-	60
	PrRd2		M		-	-	1
	PrWr2		M		-	-	1
	PrRd3		S	F	BusRd(S=1), Flush	C2	90+30
	PrWr3			M	BusUpgr	-	60
	PrRd3			M	-	-	1
	PrWr3			M	-	-	1
						TOTAL	457