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SURNAME _____

FIRST NAME _____

1) (POINTS 8/30) Consider a bus-based multicore that supports a new cache-coherence protocol called MESIF. Compared to the well-known MESI protocol, the MESIF protocol adds a 5th state called F (Forward). A copy in F state is like an S copy but the cache which holds it has the responsibility to provide (forward) the copy once a BusRd transaction involves that copy; at the same time, the S state is now simplified, as it doesn't have to respond to cache-to-cache transfers (Flush* transactions) due to a BusRd. There will be at most one cache holding a copy in the F state; on a PrRd miss, if there are other shared copies, the state will be F, whilst if another F copy observes the BusRd will go to S. Both F and S states are clean shared states. If the copy in F state is eventually evicted, there will be no copy in the F state: in this case the copy will be provided by the memory.

1a [Points 8/30] Draw the diagram of the MESIF protocol according to the above description.



1b) [Points 22/30] Assuming a cost of 1cc (1 clock-cycle) for read/write operations, 90cc for BusRd or BusRdx transactions, 60cc for BusUpgr, 20 cc for Flush* and 30cc for Flush. Evaluate the total cost (in clock-cycles) for the following streams:

stream-1 MESIF	Core Operation	C1	C2	C3	Bus Transaction	Data from	Cycles
	PrRd1						
	PrWr1						
	PrRd1						
	PrWr1						
	PrRd2						
	PrWr2						
	PrRd2						
	PrWr2						
	PrRd3						
	PrWr3						
	PrRd3						
	PrWr3						
	TOTAL						
stream-2 MESIF	Core Operation	C1	C2	C3	Bus Transaction	Data from	Cycles
	PrRd1						
	PrRd2						
	PrRd3						
	PrWr1						
	PrWr2						
	PrWr3						
	PrRd1						
	PrRd2						
	PrRd3						
	PrWr3						
	PrWr1						
	TOTAL						
stream-3 MESIF	Core Operation	C1	C2	C3	Bus Transaction	Data from	Cycles
	PrRd1						
	PrRd2						
	PrRd3						
	PrRd3						
	PrWr1						
	PrWr1						
	PrWr1						
	PrWr2						
	PrWr3						
	TOTAL						

F: state: PrRd and PrWr are exactly like for the S-state (remember that F is a shared clean state); like in S-state, the BusRdX and BusUpgr transaction cause a local invalidation, with the Flush* following the BusRdX; we enter the F-state on a read-miss once the shared-line indicates that there exist other copies in the system (S=1), so the newly acquired copy is shared; when a BusRd is snooped, this cache has the responsibility to provide the copy via a Flush* and its state will be then demoted to S (its duty is finished, the other cache with the shared copy will get the new copy in F state – it is the only possibility for a BusRd – the other is with S=0, which is not this case).