Queuing Theory

Overview of Queuing Theory Lectures

Introduction to Queueing Systems
- Important Questions about Queueing Systems
- Components of a Queueing System

Fundamental Results for General Queueing Systems
- Terminology for General Queueing Systems
- Steady State Quantities
- Fundamental Quantities of Interest
- Basic Cost Identity
- Consequences: Little’s Law and Other Useful Relationships

Properties of Poisson Arrival Processes
- Superposition and decomposition of Poisson arrivals
- Poisson arrivals see time averages
Overview of Queuing Theory Lectures, continued

Queueing Systems with Exponential Arrivals and Service
- M/M/1 queue
- M/M/S queue
- M/M/1/C queue
- Markovian queues with general state definition

Examples: security checkpoint
Safeway checkout lines
two law practices
Stanford post office
shoeshine shop

Networks of Markovian Queues
- The Equivalence Property
- Open Jackson networks
- Closed Jackson networks

Examples: car wash
Ward’s Berry Farm
roommate network

Queueing Systems with Nonexponential Distributions
- The M/G/1 Queue
- Special Cases of M/G/1: M/D/1 and M/E_k/1 queues

Examples: McDonald’s Drive Thru
homework questions

Important Questions about Queueing Systems

- What fraction of time is each server idle?
- What is the expected number of customers in the queue? in the queue plus in service?
- What is the probability distribution of the number of customers in the queue? in the queue plus in service?
- What is the expected time that each customer spends in the queue? in the queue plus in service?
- What is the probability distribution of a customer’s waiting time in the queue? in the queue plus in service?
Components of a Queueing System

Arrival Process  Queue  Service Process

Components of a Queueing System: The Arrival Process

Characteristics of Arrival Process:
- finite or infinite calling population
- bulk or individual arrivals
- interarrival time distribution
- simple or compound arrival processes
- balking or no balking
Components of a Queueing System: The Queue

- **Queue Characteristics:**
  - finite or infinite
  - queue discipline:
    - FCFS = first-come-first-served
    - LCFS = last-come-first-served
    - Priority service order
    - Random service order

Components of a Queueing System: The Service Process

- **Characteristics of Service Process:**
  - number/ configurations of servers
  - batch or single service
  - service time distribution
  - rework
Notation for Characterizing Queueing Systems

\[ \frac{\text{distribution of interarrival times}}{\text{distribution of service times}} \]

\[ \frac{\text{number of parallel servers}}{\text{maximum number of customers allowed in system}} \]

\[ \frac{\text{size of calling population}}{\text{}} \]

- \( M \) = iid exponential
- \( D \) = iid and deterministic
- \( E_k \) = Erlang with shape parameter \( k \)
- \( GI \) = iid with general distribution

Notation for Characterizing Queueing Systems: Examples

\[ M / M / 1 / \infty / \infty \]
- single server queue with Poisson arrivals and exponential service times, a.k.a. M/M/1

\[ M / M / S / \infty / \infty \]
- S server queue with Poisson arrivals and exponential service times, a.k.a. M/M/S

\[ M / M / 1 / C / \infty \]
- same as M/M/1 with finite system capacity C

\[ M / M / 1 / K / K \]
- single server queue with Poisson arrivals from a population of K potential customers; and exponential service times
Terminology for General Queueing Systems

\[ N(t) = \text{the number of customers in the system at time } t; \text{ this includes both customers that are in the queue and being served} \]

\[ P_n(t) = \text{the probability that there are exactly } n \text{ customers in the system at time } t \]

\[ = P(N(t) = n) \quad \text{N ote: } N(t) \text{ and } P_n(t) \text{ are difficult to compute for general } t. \text{ W e will be mainly interested in their values as } t \text{ becomes infinite.} \]

\[ \lambda_n = \text{mean arrival rate of entering customers when } n \text{ customers are in the system} \]

\[ \mu_n = \text{mean service rate for overall system when } n \text{ customers are in the system} \]

Steady State Quantities

\[ P_n = \text{long run probability that there will be exactly } n \text{ customers in the system} \]

\[ = \lim_{t \to \infty} P_n(t) \]

\[ \bar{\lambda} = \text{long run average arrival rate of entering customers} \]

\[ = \sum_{n=0}^{\infty} P_n \lambda_n \]

For general queues, the steady state probabilities \( P_n \) can be difficult to compute. However, if the queue is a continuous time Markov Chain, then \( P_n = \omega_n \) (if a steady-state distribution exists).
Fundamental Quantities of Interest

$L = \text{the long run average number of customers in the system}$

$L_Q = \text{the long run average number of customers waiting in queue}$

$L_S = \text{the long run average number of customers in service}$

$L = L_S + L_Q$

More Fundamental Quantities of Interest

$W' = \text{the waiting time in the system for an arbitrary customer (random variable)}$

$W'_Q = \text{the waiting time in the queue for an arbitrary customer (random variable)}$

$W = \text{the long run average amount of time a customer spends in the system}$

$= E(W')$

$W_Q = \text{the long run average amount of time a customer spends in the queue}$

$= E(W'_Q)$

$W_S = \text{the long run average amount of time a customer spends in service}$

$W = W_S + W_Q$
Basic Cost Identity

Suppose each entering customer must pay money to the system. Then the following identity applies:

\[ \text{average rate at which the system earns} = \lambda \times \text{average amount an entering customer pays} \]

Many fundamental relationships about queueing system performance can be derived from this identity. It can be used not only to compute monetary income rates but also to derive relationships between the fundamental quantities we are interested in for queueing systems. For example, if each customer pays $1 per unit time that they are in the system (i.e. either in queue or being served) then

\[ \text{the average rate at which the system earns} = L \]
\[ \text{the average amount an entering customer pays} = W \]

For this example, the basic cost identity tells us:

\[ L = \lambda W \]

Little’s Law

A quick diversion

A few topics I want to weave in before we proceed with Little’s Law and its variations:

- The superposition and decomposition of Poisson processes
- Poisson arrivals see time averages
- Computing steady state equations: rate in = rate out
Superposition of Poisson Arrival Processes

Often the arrival process to a queue consists of multiple different arrival processes of customers from different origins. It turns out that it is easy to deal with such compound arrival processes when each individual arrival process is Poisson and is independent of the others.

Suppose you have two independent Poisson arrival processes $X$ and $Y$ with respective rates $\lambda_x$ and $\lambda_y$. Then the combination of the two arrival streams is also a Poisson process with rate $\lambda_x + \lambda_y$.

Why? Let $T$ be a random variable representing the remaining time until an arrival of either type occurs. Let $T_x$ and $T_y$ be the remaining times until arrivals of type $x$ and $y$ occur. Then $T = \min(T_x, T_y)$. Since $T_x$ and $T_y$ are exponentially distributed with parameters $\lambda_x$ and $\lambda_y$ respectively, $T$ must also be exponentially distributed with parameter $\lambda_x + \lambda_y$. Hence, the combined stream has exponential interarrival times, and is thus a Poisson process with rate $\lambda_x + \lambda_y$.

Decomposition of Poisson Arrival Processes

Another nice property of Poisson arrival processes allows us to disaggregate the process into independent Poisson processes. Suppose customers arrive to our system according to a Poisson process with rate $\lambda$. Suppose that an arriving customer is of type $k$ with probability $p_k$, where

$$\sum_{k=1}^{\infty} p_k = 1.$$ 

Then the arrivals of customers of type $k$ follow a Poisson process with rate $p_k \lambda$.

An example: suppose arriving customers to a system balk with fixed probability $p$ independent of the number of people in the system. If the arrival process is Poisson with rate $\lambda$, the arrivals who stay follow a Poisson process with rate $(1-p)\lambda$. The customers who balk do so according to a Poisson process with rate $p\lambda$. 
Poisson Arrivals See Time Averages (PASTA)

The PASTA Property

For a queuing system with Poisson arrivals, an arriving customer sees the time-average steady-state number-in-system process. In other words, the long run fraction of customers that arrive to find exactly $k$ customers (not including him or herself) is given by $\omega_k$.

This property is called "Poisson Arrivals See Time Averages." It depends critically on the assumption of having a Poisson arrival process. Suppose, instead, for example, that we have deterministic arrival and service processes with respective rates $\lambda$ and $\mu$ (with $\lambda < \mu$).

The time-average steady-state fraction of time in which there are 0 customers in the system is $1 - \lambda/\mu$, and the steady-state fraction of time in which there is 1 customer is $\lambda/\mu$. But 100% of arriving customers find no customers in the system! "PASTA" does not apply in this case because the arrivals are not Poisson.

Example: Poisson Arrivals See Time Averages

**EXAMPLE: The M/ M/1 queue**

In an M/ M/1 queue, what fraction of arriving customers have to wait? The fraction that finds the system in any state $k > 0$, which, according to the PASTA property, is $1 - \omega_0 = \rho$. 

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**Poisson Arrivals See Time Averages (PASTA)**

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Computing the steady state distribution:
Rate In = Rate Out

When studying CTMCs, we learned that each of the steady state equations

\[ \sum_{i=0}^{\infty} q_{ij} \omega_i = \sum_{i=0}^{\infty} q_{ij} \omega_j = 0 \text{ for } j = 0, \ldots, S \]

could be expressed as:

\[ \text{steady state rate out of state } j = q_j \omega_j = \sum_{i=j}^{\infty} q_{ij} \omega_i = \text{steady state rate into state } j \]

In simple queuing systems that can be modeled as birth and death processes, the “rate-in = rate-out” equations are particularly simple to write down because transitions can only occur to adjacent states. In such cases, they are a shortcut to writing down the steady state equations.

Rate In = Rate Out
for the General Birth and Death Process

\[ \lambda_0 \omega_0 = \mu_1 \omega_1 \]
\[ \lambda_1 \omega_0 + \mu_2 \omega_2 = (\lambda_1 + \mu_1) \omega_1 \]
\[ \ldots \]
\[ \lambda_{n-2} \omega_{n-2} + \mu_n \omega_n = (\lambda_{n-1} + \mu_{n-1}) \omega_{n-1} \]
\[ \lambda_{n-1} \omega_{n-1} + \mu_n \omega_n = (\lambda_n + \mu_n) \omega_n \]

...
Example: Rate In = Rate Out for the M/ M/ 1 queue

\[
\begin{align*}
\text{State} & \quad \text{Rate In = Rate Out} \\
0 & \quad \mu \omega_0 = \lambda \omega_0 \\
1 & \quad \lambda \omega_0 + \mu \omega_1 = (\lambda + \mu) \omega_1 \\
\vdots & \quad \vdots \\
n-1 & \quad \lambda \omega_{n-2} + \mu \omega_{n-1} = (\lambda + \mu) \omega_{n-1} \\
n & \quad \lambda \omega_{n-1} + \mu \omega_n = (\lambda + \mu) \omega_n \\
\vdots & \quad \vdots 
\end{align*}
\]

Little’s Law

\[L = \bar{\lambda} W\]
A Proof of Little’s Law

Let $A_i$ and $D_i$ be the arrival and departure times of customer $i$. Then $W'_i = D_i - A_i$. Define the indicator function $I(A_i \leq s \leq D_i) = \begin{cases} 1 & \text{if } A_i \leq s \leq D_i, \\ 0 & \text{otherwise} \end{cases}$

\[
\frac{1}{t} \int_0^t N(s)ds = \frac{1}{t} \sum_{i=1}^{Q(t)} I(A_i \leq s \leq D_i)ds \\
= \frac{1}{t} \sum_{i=1}^{Q(t)} \int I(A_i \leq s \leq D_i)ds \\
= \frac{1}{t} \sum_{i=1}^{Q(t)} W'_i = \frac{Q(t)}{t} \sum_{i=1}^{Q(t)} W'_i \\
\lim_{t \to \infty} \frac{1}{t} \int_0^t N(s)ds = \lim_{t \to \infty} \frac{Q(t)}{t} \sum_{i=1}^{Q(t)} W'_i 
\Rightarrow L = \lambda W
\]

Other Useful Relationships

Another useful relationship can be obtained from the basic cost identity by assuming each customer pays $1 per unit time that they are in the queue. Then

- the average rate at which the system earns = $L_Q$
- the average amount an entering customer pays = $W_Q$

In this case, the basic cost identity tells us:

\[
L_Q = \lambda W_Q
\]
Other Useful Relationships

By assuming instead that each customer pays $1 per unit time that they are in service, we have
the average rate at which the system earns = \( L_s \)
the average amount an entering customer pays = \( W_s \)

The basic cost identity amounts to:

\[
L_s = \lambda W_s
\]

Exponential Queuing Models

We now show how to apply these results to queues having exponentially distributed interarrival times and service times. These types of queues are the most tractable mathematically.

Some of the queues we’ll look at:

- \( M/M/1 \)
- \( M/M/S \)
- \( M/M/1/C \) (finite capacity C)
- Queues with more general state definition
Exponential Queuing Models: The M/M/1 Queue

The simplest example of a queue that has exponentially distributed interarrival and service times is the by-now familiar M/M/1 queue. This is a queue for which the fundamental quantities are easy to compute. We will do so and also apply Little’s law to the M/M/1 case.

Define $\rho = \frac{\lambda}{\mu}$. This quantity is typically called the traffic intensity for the M/M/1 queue. We assume henceforth that $\rho < 1$, in which case we know a steady state distribution exists and is given by:

$$P_j = \omega_j = \rho^j (1 - \rho) \text{ for } j = 0, 1, \ldots$$

The arrival rate of entering customers is always $\lambda$; therefore $\overline{\lambda} = \lambda$.

Exponential Queuing Models: The M/M/1 Queue

The long run average number of people in the system can be computed using the steady state distribution:

$$L = \sum_{j=0}^{\infty} j P_j = \sum_{j=0}^{\infty} j \omega_j = (1 - \rho) \sum_{j=0}^{\infty} j \rho^j$$

$$= (1 - \rho) \frac{\rho}{(1 - \rho)}$$

$$= \frac{\rho}{(1 - \rho)}$$

$$= \frac{\lambda}{(\mu - \lambda)}$$
Exponential Queuing Models: 
The M/M/1 Queue

We can also compute the long run expected number of customers in the queue \( L_0 \) using the steady state distribution as follows:

\[
L_0 = \sum_{j=1}^{\infty} (j - 1) \omega_j = \sum_{j=1}^{\infty} j \omega_j - \sum_{j=1}^{\infty} \omega_j = \sum_{j=0}^{\infty} j \omega_j - (1 - \omega) = L - \rho = \frac{\lambda}{(\mu - \lambda)} - \frac{\lambda}{\mu} = \frac{\lambda^2}{\mu (\mu - \lambda)}
\]

The expected number of customers in service \( L_s \) is

\[
L_s = 0 \cdot \omega_0 + \sum_{j=1}^{\infty} 1 \cdot \omega_j = 1 - \omega = \rho = \frac{\lambda}{\mu}
\]

Verify that \( L = L_s + L_0 \) for this example.

Exponential Queuing Models: 
The M/M/1 Queue

The distribution for the waiting time \( W \) of an arbitrary customer (in the long run) is computed as follows:

\[
P(W \leq a) = \sum_{n=0}^{\infty} P(W \leq a | n \text{ in system when he arrives}) P(n \text{ in system when he arrives})
\]

\[
= \sum_{n=0}^{\infty} P(\text{service times of person in service} + (n-1) \text{ waiting + him \leq a}) = \sum_{n=0}^{\infty} \left[ \sum_{k=0}^{\infty} n \mu e^{-\mu} \frac{\mu^k}{n!} dt \cdot \rho^{n+1}(1-\rho) \right] = \int_0^{\infty} (\mu - \lambda) e^{-\mu} \sum_{n=0}^{\infty} \frac{(\lambda t)^n}{n!} dt = \int_0^{\infty} (\mu - \lambda) e^{-\mu} e^{\lambda t} dt = 1 - e^{-(\mu-\lambda)t} \]

So \( W \) is exponentially distributed with parameter \((\mu-\lambda)!\)
Exponential Queuing Models: The M/M/1 Queue

Now we’ll derive the distribution for the waiting time in the queue $W_Q$ of an arbitrary customer in the long run. Since an arriving customer who finds no customers in the queue has no wait, $P(W_Q = 0) = P_s = 1 - \rho$

$$P(W_Q > a) = \sum_{n=1}^{\infty} P(W_Q > a | n \text{ in system when he arrives})P(n \text{ in system when he arrives})$$

$$= \sum_{n=1}^{\infty} P(\text{service times of person in service + (n-1) waiting} > a)P_s$$

$$= \sum_{n=1}^{\infty} \int_0^\infty \mu e^{-\mu t} \left(\frac{\mu t}{(n-1)!}\right)^{n-1} dt \cdot \rho^n (1-\rho)$$

$$= \frac{\lambda (\mu - \lambda)}{\mu} e^{-\mu a} \sum_{n=1}^{\infty} \left(\frac{\lambda t}{(n-1)!}\right)^{n-1} dt$$

$$= \frac{\lambda (\mu - \lambda)}{\mu} e^{-\mu a}$$

$W_Q$ is not exponentially distributed, but the conditional distribution of $W_Q$, given that $W_Q > 0$, is!

To see this...

$$P(W_Q > a | W_Q > 0) = \frac{P(W_Q > a)}{P(W_Q > 0)} = e^{-\mu a}$$

Exponential Queuing Models: The M/M/1 Queue

Let’s now compute $W$, the expected time a customer spends in the system, first using Little’s Law:

$$W = \frac{L}{\lambda} = \frac{L_0}{\lambda} = \frac{\lambda}{\mu - \lambda} = \frac{1}{\mu - \lambda}$$

It is easy to verify that we get the same result by taking the expected value of $W$!

Using $L_0 = \bar{\lambda} W_Q$ we can compute the expected time a customer spends in the queue:

$$W_Q = \frac{L_0}{\lambda} = \frac{\bar{\lambda} \lambda^2}{\mu (\mu - \lambda)} = \frac{1}{\lambda} = \frac{\lambda}{\mu - \lambda}$$

Verify by taking the expected value of $W_Q$!

The expected time a customer spends in the service is:

$$W_s = \frac{L_s}{\lambda} = \frac{\lambda}{\mu} \cdot \frac{1}{\lambda} = \frac{1}{\mu}$$

obvious!

Evidently $W = W_s + W_Q$ for this example.
The M/M/1 Queue: Example

EXAMPLE: Security Checkpoint
Airline passengers must pass through a security checkpoint consisting of a metal detector and carry-on luggage x-ray machine. Suppose that passengers arrive according to a Poisson process with a rate of $\lambda = 10$ passengers per minute. The security checkpoint clears customers at a rate of $\mu = 12$ passengers per minute, with exponentially distributed clearance times. What is the probability that a passenger will have to wait before being checked?

$$P(W_Q > 0) = \rho e^{\rho} = \frac{5}{6}$$

(or equivalently $1 - P(W_Q = 0) = \rho$)

On average, how many customers are waiting in line to be checked?

$$L_Q = \frac{1}{\rho} = \frac{1}{12(12 - 10)} = \frac{2.5}{6}$$

On average, how long will a customer be detained at the checkpoint?

$$W = \frac{1}{\mu - \lambda} = \frac{1}{2}$$ minute

The M/M/1 Queue: Example

EXAMPLE: Security Checkpoint, continued
The airline in question wants to determine how many checkpoints to operate to minimize the costs associated with operation and customer delays. They estimate the cost of delaying a customer 1 hour to be $4. It costs $100 per hour to staff and operate a checkpoint. Assume that a passenger is equally likely to enter each checkpoint.

Let $n$ be the number of separate checkpoints. Then a given passenger goes to any particular checkpoint with probability $1/n$. We can model each of the $n$ checkpoints as a separate M/M/1 queue with a Poisson arrival process having rate $\lambda/n$. For any particular checkpoint (say, checkpoint $i$), the average number of customers present is

$$L_i = \frac{\lambda}{\mu - (\lambda/n)} = \frac{\lambda}{\mu n - \lambda} = \frac{10}{12 n - 10}$$

The expected average cost of operation plus passenger delays per hour that the airline faces when they have $n$ checkpoints open is then

$$F(n) = 100 n + 4 L = 100 n + 4 (\frac{\epsilon L}{\lambda}) = 100 n + \frac{40}{12 n - 10}$$

where $L$ represents the average total number of people at all checkpoints.
The M/M/1 Queue: Example

**EXAMPLE: Security Checkpoint**
To find the optimal number of checkpoints to operate, we need to find the minimum of the function $F(n)$.

$$F(n) = 100n + \frac{40n}{12n - 10}$$

$$F'(n) = 100 - \frac{400}{(12n - 10)^2}$$

$$F''(n) = \frac{9600}{(12n - 10)^3} \geq 0$$

Clearly the function $F(n)$ is convex, so to find its minimum we need only to set its derivative equal to zero.

$$F'(n^*) = 0 \iff (12n^* - 10)^2 = 4 \iff n^* = \pm 1$$

Since we must have at least one checkpoint, the optimal number of checkpoints is 1.

Exponential Queuing Models:
The M/M/S Queue

Another important example of a queue with exponentially distributed interarrival and service times is the M/M/S queue. For this case we'll now compute the fundamental quantities and apply Little's law.

Define $\rho = \frac{\lambda}{S\mu}$ as the traffic intensity for the M/M/S queue. To ensure the existence of a steady state distribution, we assume that $\rho < 1$. In this case:

$$P_0 = \omega_0 = (1 + K)^{-1}$$

where

$$P_j = \omega_j = \begin{cases} \frac{\lambda^j}{j!\mu^j}(1 + K)^{-1} & \text{for } j \leq S \\ \frac{\lambda^S}{S!\mu^S}(1 + K)^{-1} & \text{for } j > S \end{cases}$$

$$K = \sum_{i=1}^{S-1} \frac{\lambda^i}{i!\mu^i} + \frac{\lambda^S}{S!\mu^S} \left(1 - \frac{\lambda}{S\mu}\right)^{-1}$$

Again the arrival rate of entering customers is always $\lambda$; therefore $\bar{\lambda} = \lambda$. 
Exponential Queuing Models:
The M/ M/ S Queue

We can compute the long run expected number of customers in the queue \( L_Q \) using the steady state distribution as follows:

\[
L_Q = \sum_{j=0}^{\infty} (j - S) P_j = (1 + K)^{-1} \sum_{j=0}^{\infty} (j - S) \left( \frac{\lambda/\mu}{S!} \right)^j
\]

\[
= (1 + K)^{-1} \sum_{i=0}^{\infty} i \left( \frac{\lambda/\mu}{S!} \right)^i S^i
\]

\[
= (1 + K)^{-1} \left( \frac{\lambda/\mu}{S!} \right)^S S \sum_{i=0}^{\infty} i \rho^i
\]

\[
= (1 + K)^{-1} \left( \frac{\lambda/\mu}{S!} \right)^S \rho \frac{1}{(1 - \rho)^2}
\]

Now we can apply Little’s law and our other identities to compute:

\[
W_Q = L_Q / \lambda, \quad W = W_Q + (1/\mu), \quad L = \bar{\lambda} W
\]

M/ M/ S Queue: Example

EXAMPLE: Safeway Checkout Lines

Safeway is trying to decide how many checkout lines to keep open. An average of 18 customers arrive per hour according to a Poisson process and go to the first empty checkout line. If no checkout line is empty, suppose that arriving customers form a single line to wait for the next free checkout line. The checkout time for each customer is exponentially distributed with mean 4 minutes. It costs $20 per hour to operate a checkout line and Safeway estimates that it costs them $0.25 for each minute a customer waits in the cash register area. How many registers should the store have open?

The number \( S \) of registers open should certainly be enough to ensure that \( \lambda/Sp < 1 \); otherwise the queue will explode. This is true only if \( S > \lambda/\mu = 6/5 \). Since we can only have an integer number of servers, we require \( S \) to be at least 2.

Now we want to determine the optimal number of servers; i.e., the number that minimize Safeway’s total costs. We need to compute the expected costs incurred per hour.

First we consider the alternatives of either having 2 or 3 checkout lines open.
M/M/S Queue: Example

**case a: two checkout lines**

If there are 2 open checkout lines, then we plug $S=2$, $\lambda=18$ and $\mu=15$ into the formula for $K$:

$$K = \sum_{i=1}^{\infty} \frac{\lambda^i}{i!\mu^i} + \frac{\lambda^2}{S!\mu^S} \left(1 - \frac{\lambda}{S\mu}\right)^{-1}$$

and obtain $K = 3$. Then we can compute $L_Q$ from:

$$L_Q = (1 + K)^{-1} \left(\frac{\lambda}{\mu}\right)^{\rho} \left(1 - \rho\right)^{-1} \left(1 - \frac{\lambda}{S\mu}\right)$$

(Recall $\rho = \frac{\lambda}{S\mu}$)

which yields $L_Q = 0.675$. From this we compute

$$W_Q = L_Q / \lambda = L_Q / \lambda = 0.675 / 18 = 0.0375$$

From this and the fact that $W = W_Q + (1/\mu)$, we find that the average length of time a person waits in the register area is $W = 0.104$ hours, and the average number of people in the register area is $L = \frac{\lambda W}{\lambda} = \lambda W = 1.8738$.

Since each open checkout line costs Safeway $20 per hour, and each hour that a customer waits costs $15, the expected average costs per hour are

$$20S + 15L = 68.11$$

**case b: three checkout lines**

If instead there 3 servers, then after some arithmetic we find

$$K = 2.64$$

$$L_Q = 0.0879$$

$$W_Q = L_Q / \lambda = L_Q / \lambda = 0.0879 / 18 = 0.0049$$

$$W = W_Q + (1/\mu) = 0.0716$$

and

$$L = \frac{\lambda W}{\lambda} = \lambda W = 1.2888$$

The expected average costs per hour in the 3 server case are

$$20S + 15L = 79.33$$

The additional cost of the third checkout line exceeds the benefit of shorter lines, so two checkout lines is preferable to three. In fact, 2 is the optimal number of checkout lines. (Why?)
Exponential Queuing Models: 
The M/ M/ 1/ C Queue

Now let’s consider the variation on the M/ M/ 1 queue in which the system has capacity C. In this type of queue, customers that arrive and find C people in the system leave. The birth and death rates are:

$\lambda_i = \lambda$ for $i = 0, 1, \ldots, C - 1$
$\mu_i = \mu$ for $i = 1, 2, \ldots, C$

Once again we let $\rho = \lambda / \mu$. Since the queue is finite, a steady state distribution exists regardless of whether $\rho < 1$, and can be derived using Theorem 2 from CTMCs. If $\rho < 1$, then the steady state distribution is:

$p_j = \omega_j = \frac{1 - \rho}{1 - \rho^{C+1}} \rho^j$ for $j = 0, 1, \ldots, C$

(The case $\rho = 1$ is left as an exercise for the reader.)

The average arrival rate of entering customers (when $\rho < 1$) is then:

$\bar{\lambda} = \sum_{j=0}^{C} \lambda_j p_j = \sum_{j=0}^{C-1} \lambda p_j = \lambda \left(1 - p_C\right)$

$= \lambda \left(1 - \frac{1 - \rho}{1 - \rho^{C+1}} \rho^C\right) = \lambda \left(\frac{1 - \rho^C}{1 - \rho^{C+1}}\right)$

Exponential Queuing Models: 
The M/ M/ 1/ C Queue

The long run average number of people in the system (when $\rho < 1$) is:

$L = \sum_{j=0}^{C} j p_j = \sum_{j=0}^{C} j \omega_j = \frac{1 - \rho}{1 - \rho^{C+1}} \sum_{j=0}^{C} j \rho^j$

$= \frac{1 - \rho}{1 - \rho^{C+1}} \sum_{j=0}^{C} \rho \frac{d}{d\rho} \left(\rho^j\right)$

$= \frac{1 - \rho}{1 - \rho^{C+1}} \rho \frac{d}{d\rho} \left(\sum_{j=0}^{C} \rho^j\right)$

$= \frac{1 - \rho}{1 - \rho^{C+1}} \rho \frac{d}{d\rho} \left(\frac{1 - \rho^{C+1}}{1 - \rho}\right)$

$= \frac{\rho}{1 - \rho} - \frac{(C + 1) \rho^{C+1}}{1 - \rho^{C+1}}$
Exponential Queuing Models: The M/ M/ 1/ C Queue

The expected number of customers in the queue $L_0$ is:

$$L_0 = \sum_{j=1}^{\infty} (j-1) P_j = \sum_{j=1}^{\infty} j \omega_j - \sum_{j=1}^{\infty} \omega_j$$

$$= \sum_{j=0}^{\infty} j \omega_j - (1 - \omega_0)$$

$$= L - (1 - \omega_0)$$

The expected number of customers in service $L_s$ is

$$L_s = 0 \cdot \omega_0 + \sum_{j=1}^{\infty} 1 \cdot \omega_j = 1 - \omega_0$$

Clearly $L = L_0 + L_0$ for this example.

From Little's law and the other identities, the quantities $W$, $W_0$, and $W_s$ are easily derived.

$$W_0 = L_0 / \lambda \quad W = W_0 + (1/\mu) \quad L = \lambda W$$

The M/ M/ 1/ C Queue: Example

**EXAMPLE: Two Law Practices**

Consider two San Francisco lawyers. Lawyer 1 works with only 1 client at a time. If a second client asks for his services while he is helping the first, he will turn that client away. He charges $20,000 per client regardless of how long the client’s case takes. Lawyer 2 also helps only one client at a time, but he never turns away a client. Clients queue up to wait for his services. He charges the client he’s working with a daily rate of $300 per day.

Each lawyer receives inquiries from prospective clients according to a Poisson process with rate $\lambda = 0.01$ per day. Also, for each lawyer, the time to finish a case is exponentially distributed with mean 50 days.

Which lawyer makes more money?

We use the basic cost identity to determine the (long-run) average daily income of each lawyer:

$$\text{average rate lawyer earns} = \lambda \times \text{average amount an accepted client pays}$$
The M/ M/ 1/ C Queue: Example

**EXAMPLE: Two Law Practices**

Let $\bar{x}_i$ be the average arrival rate of clients and $F_i$ be the average daily income of lawyer $i$, $i=1,2$.

Since we know for an M/M/1/C queue

$$\lambda = \lambda \left( \frac{1 - \rho}{1 - \rho + \rho r} \right)$$

we have for lawyer 1:

$$\bar{x}_1 = \lambda (1-P_1) = \lambda \left( \frac{1 - \rho}{1 - \rho + \rho} \right) = \lambda \left( \frac{1}{1 + \rho} \right) = \left( \frac{\lambda \mu}{\mu + \lambda} \right) = \left( \begin{array}{ccc} 0.0 & 0.0 & 2 \end{array} \right)$$

$$F_1 = \bar{x}_1 \times \$20,000 = \$133.33 \text{ per day}$$

Lawyer 2's system is an M/M/1 queue:

$$F_2 = \bar{x}_2 \times \left( \frac{\$300}{\mu} \right) = \lambda \times \left( \frac{\$300}{\mu} \right) = \$150 \text{ per day}$$

Lawyer 2 makes more money.

Exponential Queuing Models with More General State Definition

In all of the examples we have seen so far, we defined the state of the system to be the number of customers in the system. This worked because in these examples, the rate of departure from the system was dependent only on the number of customers in the system. For example, in the M/M/1 server queue, if the number of people in the system is $j>S$, all servers are busy and $j-S$ people are in line. If $j<S$, then $j$ servers are busy. Since all servers are assumed to be indistinguishable, we do not care which $j$ servers are busy.

In some service systems, we may have multiple nonidentical servers. In such systems it is not sufficient to define the state as simply the number of customers in the system. The departure rates from the system depend on which specific servers are busy. We now consider two examples of exponential queues in which an extended state definition is required. We will see that all of the techniques for analyzing CTMCs and standard exponential queues still apply.
Exponential Queuing Models with More General State Definition

**EXAMPLE: Stanford Post Office**
Suppose the Stanford post office has 2 employees working at the counter, one faster than the other. Customers arrive according to a Poisson process with rate $\lambda$ and form one line. Assume there is room for at most 4 customers in the Post Office. Server 1 has service times that are exponentially distributed with rate $\mu_1$ and server 2 completes service in times that are exponentially distributed with rate $\mu_2$. When both servers are free, arriving customers choose the faster server.

How do we define the number of people in the system in this problem? If there are 2 or more customers in the system at a particular time, then we know both servers are busy. If there are no customers in the system, then no servers are busy. But if there is exactly one customer in the system, which server is busy? We need to distinguish between the two possibilities in that case.

Define states of the process to be:

- **0**: no customers present
- **(1,0)**: 1 customer present, being served by server 1
- **(0,1)**: 1 customer present, being served by server 2
- **$n$**: $n$ customers present, $n=2,3,4$

**EXAMPLE: Stanford Post Office, continued**

\[
Q = \begin{bmatrix}
-\lambda & \lambda & 0 & 0 & 0 & 0 \\
\mu_1 & -\lambda & 0 & 0 & 0 & 0 \\
\mu_2 & -\lambda & 0 & 0 & 0 & 0 \\
0 & \mu_1 & 0 & -\lambda & 0 & 0 \\
0 & 0 & \mu_1 & 0 & -\lambda & 0 \\
0 & 0 & 0 & 0 & 0 & -\lambda
\end{bmatrix}
\]
Exponential Queuing Models with More General State Definition

**EXAMPLE: Stanford Post Office, continued**

\[ 0 = -\lambda P_0 + \mu_1 P_{(1,0)} + \mu_2 P_{(0,1)} \]
\[ 0 = \lambda P_0 - (\lambda + \mu_1) P_{(1,0)} + \mu_1 P_{(0,1)} \]
\[ 0 = \lambda P_{(1,0)} - (\lambda + \mu_1) P_{(2,0)} + \mu_1 P_{(0,1)} \]
\[ 0 = \lambda P_{(2,0)} - (\lambda + \mu_1) P_{(3,0)} + (\mu_1 + \mu_2) P_{(0,1)} \]
\[ 0 = \lambda P_{(3,0)} - (\mu_1 + \mu_2) P_{(0,1)} \]

1 = \mathbf{P}_0 + P_{(1,0)} P_{(0,1)} + P_2 + P_3 + P_4

Suppose \( \lambda = \mu_2 = 1 \) per minute and \( \mu_1 = 2 \) per minute. In this case, solving with Matlab gives us:

\(( P_0, P_{(1,0)}, P_{(0,1)}, P_2, P_3, P_4 ) = ( 0.5294, 0.2118, 0.1059, 0.1059, 0.0353, 0.0118 )\)

---

**Exponential Queuing Models with More General State Definition**

**EXAMPLE: Stanford Post Office, continued**

From the steady state distribution, we can answer a number of interesting questions:

1. What proportion of time is each server busy?

   - Server 1’s proportion of busy time is \( P_{(1,0)} + P_2 + P_3 + P_4 = 0.3647 \)
   - Server 2’s proportion of busy time is \( P_{(0,1)} + P_2 + P_3 + P_4 = 0.2588 \)

2. What is the average number of customers in service at the Post Office?

   \[ L = 0 \cdot P_0 + 1 \cdot P_{(1,0)} + P_{(0,1)} + 2 \cdot ( P_2 + P_3 + P_4 ) = 0.6824 \]

3. What is the average time a customer spends in service at the Post Office?

   \[ \bar{X} = \lambda ( P_0 + P_{(1,0)} + P_{(0,1)} + P_2 + P_3 ) = 1 \cdot P_4 = 0.8941 \]
   \[ W = L / \bar{X} = 0.6824 / 0.8941 = 0.7632 \text{ minutes} \]
Exponential Queuing Models with More General State Definition

**EXAMPLE: Shoeshine Shop**
Consider a shoeshine shop with 2 chairs. An entering customer goes to chair 1 to get his shoes polished. When the polishing is done, he will either go on to chair 2 to have his shoes buffed if that chair is empty, or else wait in chair 1 until chair 2 becomes empty. A potential customer will enter the shop only if chair 1 is empty. Potential customers arrive according to a Poisson process with rate $\lambda$. The service time in chair $i$ exponentially distributed with rate $\mu_i$ for $i=1,2$.

We'd like to know:

1. What proportion of potential customers enter the shop?
2. What is the mean number of customers in the shop?
3. What is the average amount of time an entering customer spends in the shop?

The state of the system must include more information than simply the number of customers in the shop, because the arrival and departure rates depend on where in the shop the customers are.

---

**EXAMPLE: Shoeshine Shop, continued**

Define states of the process to be:

- $0$: no customers present
- $(1,0)$: 1 customer present, in chair 1
- $(0,1)$: 1 customer present, in chair 2
- $(1,1)$: two customers present, both being served
- $(b,1)$: two customers present, but the customer in chair 1 is waiting for chair 2 to become available

![State Transition Diagram](image)
Exponential Queuing Models with More General State Definition

**EXAMPLE: Shoeshine Shop, continued**

The transition rate matrix for this process is

\[
Q = \begin{pmatrix}
-\lambda & \lambda & & \\
\mu_1 & -\mu_1 & \mu_2 & \\
& \mu_2 & -\mu_2 & \lambda \\
& & \mu_2 & -\mu_2 \\
\end{pmatrix}
\]

which yields the steady state equations:

\[
\begin{align*}
0 &= -\lambda P_0 + \mu_1 P_{(0,1)} \\
0 &= \lambda P_0 - \mu_1 P_{(1,0)} + \mu_2 P_{(1,1)} \\
0 &= \mu_1 P_{(1,0)} - (\lambda + \mu_1) P_{(0,1)} + \mu_2 P_{(0,1)} \\
0 &= \lambda P_{(1,0)} - (\mu_1 + \mu_2) P_{(1,1)} \\
0 &= \mu_1 P_{(1,1)} - \mu_2 P_{(2,1)} \\
P_0 + P_{(0,1)} + P_{(1,0)} + P_{(1,1)} + P_{(2,1)} &= 1
\end{align*}
\]

Exponential Queuing Models with More General State Definition

**EXAMPLE: Shoeshine Shop, continued**

We are particularly interested in what proportion of potential customers enter the shop:

\[
P_0 + P_{(0,1)}
\]

the average number of customers in the shop:

\[
L = P_{(0,1)} + P_{(1,0)} + 2(P_{(1,1)} + P_{(2,1)})
\]

as well as the average time a customer spends in the shop:

\[
W = \frac{L}{\lambda} = \frac{P_{(0,2)} + P_{(1,0)} + 2(P_{(1,2)} + P_{(2,1)})}{\lambda (P_0 + P_{(0,1)})}
\]
Exponential Queuing Models with More General State Definition

**EXAMPLE: Shoeshine Shop, continued**

Let's suppose that the arrival rate $\lambda = 1$ customer per minute. Consider first the case that $\mu_1 = 1$ and $\mu_2 = 2$, i.e., the server for chair 2 is twice as fast. In that case, the steady state equations are:

\[
\begin{align*}
P_0 &= 2P_{(0,1)} \\
P_{(1,0)} &= P_0 + 2P_{(1,2)} \\
3P_{(0,1)} &= P_{(1,0)} + 2P_{(0,1)} \\
P_{(0,1)} &= 3P_{(1,1)} \\
P_{(1,1)} &= 2P_{(2,1)} \\
P_0 + P_{(0,1)} + P_{(1,0)} + P_{(1,1)} + P_{(0,2)} &= 1
\end{align*}
\]

Solution:

\[
\begin{align*}
P_0 &= \frac{12}{37} \\
P_{(1,0)} &= \frac{16}{37} \\
P_{(0,1)} &= \frac{6}{37} \\
P_{(1,1)} &= \frac{2}{37} \\
P_{(0,2)} &= \frac{1}{37}
\end{align*}
\]

From the steady state distribution, we can answer our questions:

1. the proportion of potential customers entering the shop = $P_0 + P_{(0,1)} = \frac{18}{37}$

2. $L = \frac{22}{37} + 2\left(\frac{3}{37}\right) = \frac{28}{37}$

3. $W = \frac{L}{\lambda(P_0 + P_{(0,1)})} = \frac{28}{37} \cdot \frac{37}{18} = \frac{14}{9}$

Exponential Queuing Models with More General State Definition

**EXAMPLE: Shoeshine Shop, continued**

Now suppose that $\mu_1 = 2$ and $\mu_2 = 1$, i.e., the server for chair 1 is twice as fast. In that case, the steady state equations are:

\[
\begin{align*}
P_0 &= P_{(0,1)} \\
2P_{(1,0)} &= P_0 + P_{(1,2)} \\
2P_{(0,1)} &= 2P_{(1,0)} + P_{(0,1)} \\
P_{(0,1)} &= 3P_{(1,1)} \\
2P_{(1,1)} &= P_{(2,1)} \\
P_0 + P_{(0,1)} + P_{(1,0)} + P_{(1,1)} + P_{(0,2)} &= 1
\end{align*}
\]

Solution:

\[
\begin{align*}
P_0 &= \frac{3}{11} \\
P_{(1,0)} &= \frac{2}{11} \\
P_{(0,1)} &= \frac{3}{11} \\
P_{(1,1)} &= \frac{1}{11} \\
P_{(0,2)} &= \frac{2}{11}
\end{align*}
\]

(1) the proportion of potential customers entering the shop = $P_0 + P_{(0,1)} = \frac{6}{11}$

(2) $L = \frac{5}{11} + 2\left(\frac{3}{11}\right) = 1$

(3) $W = \frac{L}{\lambda(P_0 + P_{(0,1)})} = \frac{1}{11} \cdot \frac{11}{6} = \frac{11}{6}$

Having the server for the second chair server be faster leads to loss of more potential customers, but shorter average waits and fewer customers in the shop on average.
Networks of Markovian Queues

Queueing systems often provide more than a single service. Many real-world systems involve customers receiving a number of different services, each at a different service station, where each station has a queue for service. We can model systems of this kind using networks of queues. There are two types of queueing networks: open and closed. An open queueing network is one in which customers can enter and leave the system. A closed queueing network is one in which there are a fixed number of customers that never leave; new customers never enter.

We will be concerned with networks of Markovian queues called Jackson networks. These networks have exponentially distributed service times at each service station and exponentially distributed interarrival times of new customers. The sequence of stations visited is governed by a probability matrix. We’ll consider first open and then closed Jackson networks. We’ll start off with an important result that will help with the analysis.

The Equivalence Property

The Equivalence Property

In an M/M/S queue which is positive recurrent (i.e., λ < Sμ), the steady state output of the system is a Poisson process with the same rate as the input process. Note: S may be infinite.
Open Jackson Networks: Example

Example: Car Wash (a tandem network)

From the equivalence principle, the departure process of the wash station is a Poisson process with rate $\lambda = 30$/hr in steady state. Thus, in steady state both stations are effectively $M/M/1$ queues with arrival rate $\lambda = 30$/hr.

Open Jackson Networks: Example

Example: Car Wash, continued

In order to analyze this system, we need to keep track of the number of cars at each station. Thus our state definition is a vector with 2 elements $X(t) = (n_1, n_2)$, where $n_1$ represents the number of cars being washed or in queue to be washed and $n_2$ represents the number of cars being waxed or in queue to be waxed.

Even for such a simple network, the rate diagram is a mess.

We would like to find the steady state distribution $\omega = (\omega_{n_1, n_2})$ for $X(t) = (n_1, n_2)$. To do this we would need to solve the system $\omega Q = 0$, $\omega e = 1$. But the $Q$ matrix is difficult to write down. To make our lives easier, recall that the equations $\omega Q = 0$ simply state that the rate into any state equals the rate out of that state.
Open Jackson Networks: Example

Example: Car Wash, continued

The equations $\omega Q = 0$ are given by

$$
\begin{align*}
(0,0) & : \lambda \omega_{0,0} = \\
(n_1,0) : (n_1 > 0) & : (\lambda + \mu_1) \omega_{n_1,0} = \\
(0,n_2) : (n_2 > 0) & : (\lambda + \mu_2) \omega_{0,n_2} = \\
(n_1,n_2) : (n_1,n_2 > 0) & : (\lambda + \mu_1 + \mu_2) \omega_{n_1,n_2} = \\
\end{align*}
$$

To compute the steady state probabilities $P_{n_1,n_2} = \omega_{n_1,n_2}$ we must solve the above equations along with

$$
\sum_{n_1=0}^{\infty} \sum_{n_2=0}^{\infty} P_{n_1,n_2} = \sum_{n_1=0}^{\infty} \sum_{n_2=0}^{\infty} \omega_{n_1,n_2} = 1
$$

Example: Car Wash, continued

Rather than solve the steady state equations directly, let's guess at a solution.

First, by the equivalence principal, we know station $i$ is an $M/M/1$ queue with arrival rate $\lambda$ and service rate $\mu_i$. Assuming $\rho_i = \lambda/\mu_i < 1$ for $i=1,2$, we know

$$
P(n_i \text{ at wash station}) = \left( \frac{\lambda}{\mu_i} \right)^{n_i} \left( 1 - \frac{\lambda}{\mu_i} \right) = \rho_i^{n_i} (1 - \rho_i)
$$

$$
P(n_2 \text{ at wax station}) = \left( \frac{\lambda}{\mu_2} \right)^{n_2} \left( 1 - \frac{\lambda}{\mu_2} \right) = \rho_2^{n_2} (1 - \rho_2)
$$

Now, if it were true that the number of cars at the wash and wax stations at any time were independent random variables, then we would know that

$$
P_{n_1,n_2} = P(n_1 \text{ at wash station and } n_2 \text{ at wax station})
\quad = P(n_1 \text{ at wash station}) P(n_2 \text{ at wax station}) = \rho_i^{n_1} (1 - \rho_i) \rho_2^{n_2} (1 - \rho_2)
$$
Example: Car Wash, continued

As it turns out, our guess is correct. Verify for yourself that \( \omega_{n_1,n_2} = P_{n_1,n_2} \) given by

\[
P_{n_1,n_2} = \rho_1^{n_1} (1 - \rho_1) \rho_2^{n_2} (1 - \rho_2)
\]

satisfies the steady state equations

\[
\begin{align*}
\text{state} & \quad \text{rate out} & \quad = & \quad \text{rate out} \\
(0,0) & \quad \lambda P_{0,0} & = & \quad \mu_1 P_{0,1} \\
(n_1,0); \ n_1 > 0 & \quad (\lambda + \mu_1) P_{n_1,0} & = & \quad \mu_2 P_{n_1,1} + \lambda P_{n_1,1} \\
(0,n_2); \ n_2 > 0 & \quad (\lambda + \mu_2) P_{0,n_2} & = & \quad \mu_2 P_{0,n_2} + 1 + \mu_1 P_{1,n_2} \\
(n_1,n_2); \ n_1, n_2 > 0 & \quad (\lambda + \mu_1 + \mu_2) P_{n_1,n_2} & = & \quad \mu_2 P_{n_1,n_2} + 1 + \mu_1 P_{n_1+1,n_2} + \lambda P_{n_1+1,n_2}
\end{align*}
\]

and

\[
\sum_{n_1=0}^{\infty} \sum_{n_2=0}^{\infty} P_{n_1,n_2} = 1
\]

Open Jackson Networks: Example

Example: Car Wash, continued

The result

\[
P_{n_1,n_2} = \rho_1^{n_1} (1 - \rho_1) \rho_2^{n_2} (1 - \rho_2)
\]

is called the product form solution for the steady state distribution of the state of the queueing network.

What is its significance? It tells us that the numbers of cars at the wash and wax stations at any given time are, in fact, independent random variables!

This result depends heavily on the fact that the queue in front of each station has infinite capacity. If not, the numbers of cars at the two stations would not be independent.
Open Jackson Networks: Example

**Example: Car Wash, continued**

The product form solution of the steady state distribution makes it easy to calculate the expected number of cars in the system:

\[
L = \sum_{n_1=0}^{\infty} \sum_{n_2=0}^{\infty} (n_1 + n_2) p_{n_1,n_2} = \sum_{n_1=0}^{\infty} \sum_{n_2=0}^{\infty} (n_1 + n_2) \rho_1^{n_1} (1 - \rho_1) \rho_2^{n_2} (1 - \rho_2)
\]

\[
= \sum_{n_1=0}^{\infty} n_1 \rho_1^{n_1} (1 - \rho_1) \sum_{n_2=0}^{\infty} \rho_2^{n_2} (1 - \rho_2) + \sum_{n_2=0}^{\infty} n_2 \rho_2^{n_2} (1 - \rho_2) \sum_{n_1=0}^{\infty} \rho_1^{n_1} (1 - \rho_1)
\]

\[
= \sum_{n_1=0}^{\infty} n_1 \rho_1^{n_1} (1 - \rho_1) + \sum_{n_2=0}^{\infty} n_2 \rho_2^{n_2} (1 - \rho_2)
\]

\[
= (1 - \rho_1) \frac{\rho_1}{(1 - \rho_1)^2} + (1 - \rho_2) \frac{\rho_2}{(1 - \rho_2)^2}
\]

\[
= \frac{\rho_1}{1 - \rho_1} + \frac{\rho_2}{1 - \rho_2}
\]

\[
= \frac{\lambda}{\mu_1 - \lambda} + \frac{\lambda}{\mu_2 - \lambda}
\]

Notice that this equals the sum of the expected numbers of cars at the two stations, respectively.

**Open Jackson Networks: Example**

**Example: Car Wash, continued**

Using this last result and Little’s Law, we can compute the average time a car spends at the car wash:

\[
W = \frac{L}{\lambda} = \frac{1}{\mu_1} + \frac{1}{\mu_2} = \frac{1}{10} + \frac{1}{15} = \frac{1}{6} \text{ hour} = 6 \text{ min}
\]

Notice that this equals the sum of the average time a car spends at each of the two stations, respectively.

What made the computation of \(L\) and \(W\) easy in this example was the product form solution of the steady state distribution. It turns out that there is a much more general framework under which a product form solutions are available. These are called Jackson networks.
Open Jackson Networks

An Open Jackson Network is a network of \( m \) service stations, where station \( i \) has

1. an infinite queue
2. customers arriving from outside the system according to a Poisson process with rate \( a_i \)
3. \( s_i \) identical servers, each with an exponential service time distribution with rate \( \mu_i \)
4. the probability that a customer exiting station \( i \) goes to station \( j \) is \( p_{ij} \)
5. a customer exiting station \( i \) departs the system with probability \( 1 - \sum_{j \neq i} p_{ij} \)

Open Jackson Networks: Example

**Example: Ward's Berry Farm**

1. Strawberry Field (self-service)
   - \( a_1 = 10/\text{hr} \)
   - \( \mu_1 = 2/\text{hr} \)
   - \( p_{21} = 0.1 \)

2. Berry Weighing Station
   - 1 server
   - \( \mu_2 = 60/\text{hr} \)
   - \( p_{24} = 0.6 \)

3. Farmstand (self-service)
   - \( \mu_3 = 6/\text{hr} \)
   - \( p_{34} = 0.9 \)

4. Cashiers
   - 3 servers
   - \( \mu_4 = 20/\text{hr} \)

Recall from homework: self-service is effectively equivalent to having infinitely many servers.
Open Jackson Networks: Traffic Equations

Define:

\[ \lambda_i = \text{total arrival rate to station } i \]  
(including external and internal arrivals)

The values \( \lambda_i \), \( i = 1, 2, \ldots, m \) must satisfy the following equations:

\[ \lambda_i = a_i + \sum_{j=1}^{m} \lambda_j p_{ij} \quad \text{for } i = 1, 2, \ldots, m \]

Traffic Equations for Open Jackson Networks

\[ \lambda_i = a_i + \sum_{j=1}^{m} \lambda_j p_{ij} \quad \text{for } i = 1, 2, \ldots, m \]

Open Jackson Networks: Example

Example: Ward’s Berry Farm, continued

In this example, our traffic equations look like:

\[ \lambda_1 = a_1 + \sum_{j=1}^{m} \lambda_j p_{1j} = 10 + (0.1)\lambda_2 \]
\[ \lambda_2 = a_2 + \sum_{j=1}^{m} \lambda_j p_{2j} = \lambda_1 \]
\[ \lambda_3 = a_3 + \sum_{j=1}^{m} \lambda_j p_{3j} = 30 + (0.3)\lambda_2 \]
\[ \lambda_4 = a_4 + \sum_{j=1}^{m} \lambda_j p_{4j} = (0.6)\lambda_2 + (0.9)\lambda_3 \]

Solving these simultaneously gives us:

\[ \lambda_1 = \lambda_2 = \frac{100}{9} \]
\[ \lambda_3 = 30 + \frac{3}{10} \frac{100}{9} = \frac{100}{3} \]
\[ \lambda_4 = \frac{6}{10} \frac{100}{9} + \frac{9}{10} \frac{100}{3} = \frac{110}{3} \]
Open Jackson Networks: The Product Form Solutions

The state of the system is a vector of length $m$:

$$X(t) = (n_1, n_2, n_3, ..., n_m)$$

means

- $n_1$ customers at station 1,
- $n_2$ customers at station 2,
- $n_3$ customers at station 3,
- $n_m$ customers at station $m$.

It turns out that if $\lambda_i / s_i \mu_i < 1$ for each station $i = 1, 2, ..., m$, then, as in the car wash example, the steady state distribution has a product form solution:

**Product Form Solution for Open Jackson Networks**

$$P_{n_1, n_2, ..., n_m} = P(n_1 \text{ at station 1}, n_2 \text{ at station 2}, ..., n_m \text{ at station } m)$$

$$= P(n_1 \text{ at station 1}) P(n_2 \text{ at station 2}) \cdots P(n_m \text{ at station } m)$$

Open Jackson Networks: Example

**Example: Ward’s Berry Farm, continued**

To determine the steady state distribution for the entire network in this example, we first recall what we know about the $M/M/1$, $M/M/3$, and $M/M/\infty$ queues to determine whether each of the 4 stations has a steady state distribution of its own:

- stations 1 and 3 have steady state distributions because the $M/M/\infty$ queue always does!
- station 2: $\lambda_2 = (100/9) < \mu_2 = 60 \iff$ station 2 has a steady state distribution
- station 4: $\lambda_4 = (110/3) < s_4 \mu_4 = 3(20) = 60 \iff$ station 4 has a steady state distribution
Open Jackson Networks: Example

Example: Ward’s Berry Farm, continued

We compute each station’s own steady state distribution using what we already know about the M/M/1, M/M/3 and M/M/∞ queues:

\[
P(j \text{ at station } 1) = \frac{(\lambda_i / \mu_j)^i}{i!} e^{-\lambda_i / \mu_j},
\]

\[
P(j \text{ at station } 2) = (1 - \lambda_i / \mu_2)(\lambda_i / \mu_2)^i
\]

\[
P(j \text{ at station } 3) = \frac{(\lambda_i / \mu_j)^i}{i!} e^{-\lambda_i / \mu_j},
\]

\[
P(j \text{ at station } 4) = \begin{cases} 
\frac{(\lambda_i / \mu_4)^i}{i!} (1 + K)^{-1} & \text{for } 0 \leq j \leq 3 \\
\frac{(\lambda_i / \mu_4)^i}{3!} (1 + K)^{-1} & \text{for } j \geq 4
\end{cases}
\]

where \( K = \sum_{i=0}^{3} \frac{\lambda_i / \mu_4)^i}{i!} + \frac{(\lambda_i / \mu_4)^3}{3!} \left( \frac{1}{1 - (\lambda_i / 3 \mu_4)} \right) \)

Open Jackson Networks: Example

Example: Ward’s Berry Farm, continued

We express the steady state distribution of the entire network in product form:

\[
P_{n_1, n_2, n_3, n_4} = P(n_1 \text{ at station } 1) P(n_2 \text{ at station } 2) P(n_3 \text{ at station } 3) P(n_4 \text{ at station } 4)
\]

\[
= \frac{(\lambda_i / \mu_1)^{n_1}}{n_1!} e^{-\lambda_i / \mu_1} \cdot (1 - \lambda_i / \mu_2)(\lambda_i / \mu_2)^{n_2} \\
\frac{(\lambda_i / \mu_3)^{n_3}}{n_3!} e^{-\lambda_i / \mu_3} \cdot \begin{cases} 
\frac{(\lambda_i / \mu_4)^{n_4}}{n_4!} (1 + K)^{-1} & \text{if } 0 \leq n_4 \leq 3 \\
\frac{(\lambda_i / \mu_4)^{n_4}}{3!} (1 + K)^{-1} & \text{if } n_4 \geq 4
\end{cases}
\]

This looks cumbersome but is easily evaluated for any particular state vector \((n_1, n_2, n_3, n_4)\). An example is \((n_1, n_2, n_3, n_4) = (0,0,0,0)\):

\[
P_{0,0,0,0} = e^{-\lambda_i / \mu_1} (1 - \lambda_i / \mu_2)(1 + K)^{-1}
\]

\[
= e^{-5.556} (1 - 0.185) e^{-5.556} (1 + 7.11)^{-1}
\]

\[
= 1.4962 \times 10^{-5}
\]
Open Jackson Networks

Steps to Analyze an Open Jackson Network

1. Solve the traffic equations \( \lambda_i = a_i + \sum_{j=1}^{m} \lambda_j p_{ij} \) for \( i = 1, 2, \ldots, m \).

2. Check that \( \lambda_i < s_i \mu_i \) for each station \( i = 1, \ldots, m \).
   - If \textbf{NO}, the number of customers in the network blows up, so there is no steady state distribution.
   - If \textbf{YES}, go to step 3.

3. For each station \( i \), calculate the steady state distribution \( \omega^i = (\omega^i_1, \omega^i_2, \omega^i_3, \ldots) \) for an M/M/s\(_i\) queue with arrival rate \( \lambda_i \) and the service rate \( \mu_i \).

4. The steady state probability that there are \( n_1 \) customers at station 1, \( n_2 \) customers at station 2, \ldots, \( n_m \) customers at station \( m \) is just

   \[
   P_{n_1, n_2, \ldots, n_m} = \left( \prod_{i=1}^{m} \frac{\omega^i_{n_i}}{G^i_{m_i}} \right) \left( \prod_{j=1}^{m} \frac{\rho^j_{n_j}}{G^j_{n_j}} \right) = \prod_{i=1}^{m} \frac{\omega^i_{n_i}}{G^i_{m_i}} \left( \prod_{j=1}^{m} \frac{\rho^j_{n_j}}{G^j_{n_j}} \right) 
   \]

Open Jackson Networks

In the special case that all \( n \) stations have a single server (\( s_i = 1 \) for all \( i \)) then the analysis is particularly easy:

1. Solve the traffic equations \( \lambda_i = a_i + \sum_{j=1}^{m} \lambda_j p_{ij} \) for \( i = 1, 2, \ldots, m \).

2. Check that \( \rho_i = \lambda_i / \mu_i < 1 \) for each station \( i = 1, \ldots, m \).
   - If \textbf{NO}, there is no steady state distribution.
   - If \textbf{YES}, go to step 3.

3. For each station \( i \), the steady state distribution \( \omega^i = (\omega^i_1, \omega^i_2, \omega^i_3, \ldots) \) for an M/M/1 queue with arrival rate \( \lambda_i \) and the service rate \( \mu_i \) is:

   \[
   \omega^i_j = (1 - \rho_i)(\rho_i)^j 
   \]

4. The steady state probability that there are \( n_1 \) customers at station 1, \( n_2 \) customers at station 2, \ldots, \( n_m \) customers at station \( m \) is just

   \[
   P_{n_1, n_2, \ldots, n_m} = \omega^1_{n_1} \cdot \omega^2_{n_2} \cdots \omega^m_{n_m} = \left( \prod_{i=1}^{m} (1 - \rho_i)(\rho_i)^{n_i} \right) = \prod_{i=1}^{m} (1 - \rho_i)(\rho_i)^{n_i} 
   \]
Example: Ward’s Berry Farm, continued

Ward’s Berry Farm makes, on average, $15 per hour that a single customer spends picking berries. What is their average hourly income from the berry field? It’s

\[ $15 \cdot L_i \]

where

\[ L_i = \frac{\lambda_i}{\mu_i} = \frac{1009}{2} = 5.56 \]

Thus, their average hourly income from the berry field alone is $83.40.

Open Jackson Networks: Example

Example: Ward’s Berry Farm, continued

What is the average number of people at Ward’s Berry Farm at a given point in time in steady state? By the product form solution for Open Jackson Networks, we know the answer is

\[ L = L_1 + L_2 + L_3 + L_4 \]

where \( L_i \) is the average number of people at station \( i \) in steady state.

\[ L_1 = \frac{\lambda_i}{\mu_i} = \frac{1009}{2} = 5.56 \]

\[ L_2 = \frac{\lambda_2}{(\mu_2 - \lambda_2)} = \frac{100}{(60 - (100/9))} = 0.28 \]

\[ L_3 = \frac{\lambda_3}{\mu_3} = \frac{1003}{6} = 5.56 \]

\[ L_4 = (1 + K_4)^{-1} \cdot \frac{\lambda_4 / \mu_4}{3!} \cdot \frac{\rho_4}{1 - \rho_4} = 0.20 \]

where

\[ \rho_4 = \frac{\lambda_4}{5 \mu_4} = \frac{\lambda_4}{3 \mu_4} \]

\[ L = L_1 + L_2 + L_3 + L_4 = 11.6 \]

is the average steady state number of people at the farm.
Open Jackson Networks: Example

Example: Ward’s Berry Farm, continued
What is the average duration time a visiting customer spends at the farm? It turns out we can’t just add average times they spend at each station, because the customer doesn’t necessarily visit each station exactly once. Instead, we can just use Little’s Law:

\[ L = \lambda W \]

where \( \lambda \) represents the average total external arrival rate to the entire system, i.e.,

\[ \lambda = a_1 + a_2 + a_3 + a_4 = 10 + 30 = 40 \]

So a visiting customer spends, on average,

\[ W = L / \lambda = 11.6 / 40 = .29 \] hours at the farm.

Closed Jackson Networks

A closed Jackson network of queues is a network with a fixed number \( n \) of customers and \( m \) service stations, where station \( i \) has

1. an infinite queue
2. \( s_i \) identical servers, each with an exponential service time distribution with rate \( \mu_i \)
3. the probability that a customer exiting station \( i \) goes to station \( j \) is \( p_{ij} \)
4. no customers exiting the system, i.e., \( \sum_{j=1}^{m} p_{ij} = 1 \)
5. no arrivals of customers from outside the system
Closed Jackson Networks

As in the open network case, we let

\[ \lambda_i = \text{total arrival rate to station } i \]

where this time there are no external arrivals.

Now the values \( \{\lambda_i, i=1,2,...,m\} \) must satisfy:

**Traffic Equations for Closed Jackson Networks**

\[
\begin{align*}
\lambda_i &= \sum_{j=1}^{m} \lambda_j p_{ji} \text{ for } i = 1,2,...,m \\
\sum_{i=1}^{m} \lambda_i &= 1
\end{align*}
\]

Closed Jackson Networks

**Steps to Analyze a Closed Jackson Network**

1. Solve the traffic equations \( \lambda_i = \sum_{j=1}^{m} \lambda_j p_{ji} \text{ for } i = 1,2,...,m \), \( \sum_{i=1}^{m} \lambda_i = 1 \)

   If a solution exists, go to step 2.

   If not, there is no steady state distribution.

2. For each station \( i \), calculate the steady state distribution \( \omega^i = (\omega_{s1}^i, \omega_{s2}^i, \omega_{s3}^i,...) \)

   for an M/M/1 queue with arrival rate \( \lambda_i \) and the service rate \( \mu_i \).

   (If \( \lambda_i \) exceeds \( \mu_i \) and \( s_i=1 \), then it's okay to use the steady state distribution anyway.)

continued....
Closed Jackson Networks

Steps to Analyze a Closed Jackson Network

(3) The steady state probability that there are \( n_1 \) customers at station 1, \( n_2 \) customers at station 2, \ldots, \( n_m \) customers at station \( m \) is just

\[
P_{n_1, n_2, \ldots, n_m} = \begin{cases} 0 & \text{if } n_1 + n_2 + \ldots + n_m \neq n \\ \sum_{j_1, j_2, \ldots, j_m \geq 0} \omega_{j_1}^{1} \cdot \omega_{j_2}^{2} \cdots \omega_{j_m}^{m} & \text{if } n_1 + n_2 + \ldots + n_m = n \end{cases}
\]

Closed Jackson Networks: Example

Example: Roommate Network
Suppose that between yourself and your two roommates, you own 2 Gameboys. When you get your hands on one, you tend to use it for a time that's exponentially distributed with mean of 1 hour. When you're tired of it, you give it to your roommate Barbara with probability 0.8 or your other roommate Joyce with probability 0.2. The lengths of time that Barbara and Joyce keep a Gameboy is exponentially distributed with mean of 1.5 hours and 2 hours, respectively. Barbara will always give the one she's been playing with to you when she's through, whereas Joyce give it to you only 60% of the time; the rest of the time she'll let Barbara have it next. What's the probability that each of your roommates has Gameboy, but you don't?
Closed Jackson Networks: Example

Example: Roommate network, continued
(1) First we solve the traffic equations to obtain the arrival rates are each “station”:

\[ \lambda_i = \sum_{j=1}^{m} \lambda_j p_{ji} \text{ for } i = 1, 2, \ldots, m. \]

\[ \sum_{i=1}^{m} \lambda_i = 1 \]

\[ \lambda_1 = \lambda_2 + (0.6) \lambda_3 \]
\[ \lambda_2 = (0.8) \lambda_1 + (0.4) \lambda_3 \]
\[ \lambda_3 = (0.2) \lambda_1 \]
\[ \lambda_1 + \lambda_2 + \lambda_3 = 1 \]

The resulting arrival rates are:
\[ \lambda_1 = \frac{25}{52}, \lambda_2 = \frac{22}{52}, \lambda_3 = \frac{5}{52} \]

(2) We now calculate the steady state distribution for each station i by modeling it as an M/M/1 queue with arrival rate \( \lambda_i \) and the service rate \( \mu_i \).

station 1 (you): \( \omega_1 = \left( \frac{27}{52} \right)^i \left( \frac{25}{52} \right) \)

station 2 (Barbara): \( \omega_2 = \left( \frac{19}{52} \right)^i \left( \frac{33}{52} \right) \)

station 3 (Joyce): \( \omega_3 = \left( \frac{42}{52} \right)^i \left( \frac{10}{52} \right) \)

(3) Now we want to compute the steady state probability that each of your roommates has a Gameboy. This probability is denoted by \( P_{0,1,1} \).
Closed Jackson Networks: Example

Example: Roommate network, continued

The steady state probability that there are $n_1$ Gameboys with you, $n_2$ Gameboys with Barbara, and $n_3$ Gameboys with Joyce is

$$P_{n_1,n_2,n_3} = \begin{cases} 
\omega_1 \cdot \omega_2 \cdot \omega_3 & \text{if } n_1 + n_2 + n_3 \neq 2 \\
\sum_{j_1+j_2+j_3=2} \omega_1^{j_1} \cdot \omega_2^{j_2} \cdot \omega_3^{j_3} & \text{if } n_1 + n_2 + n_3 = 2
\end{cases}$$

Let's first compute the denominator:

$$\sum_{j_1+j_2+j_3=2} \omega_1^{j_1} \cdot \omega_2^{j_2} \cdot \omega_3^{j_3}$$

Closed Jackson Networks: Example

Example: Roommate network, continued

There are six feasible combinations of indices $(j_1, j_2, j_3)$ that appear in the sum, namely those satisfying $j_1, j_2, j_3 \geq 0$, $j_1 + j_2 + j_3 = 2$

<table>
<thead>
<tr>
<th>$(j_1, j_2, j_3)$</th>
<th>$\omega_1^{j_1} \cdot \omega_2^{j_2} \cdot \omega_3^{j_3}$</th>
</tr>
</thead>
<tbody>
<tr>
<td>$(2, 0, 0)$</td>
<td>$\omega_1^2 \cdot \omega_3^0 = \left(\frac{27}{52}\right)^2 \left(\frac{52}{52}\right)^0$</td>
</tr>
<tr>
<td>$(0, 2, 0)$</td>
<td>$\omega_1^0 \cdot \omega_2^2 \cdot \omega_3^0 = \left(\frac{27}{52}\right)^0 \left(\frac{19}{52}\right)^2 \left(\frac{52}{52}\right)^0$</td>
</tr>
<tr>
<td>$(0, 0, 2)$</td>
<td>$\omega_1^0 \cdot \omega_2^0 \cdot \omega_3^2 = \left(\frac{27}{52}\right)^0 \left(\frac{19}{52}\right)^0 \left(\frac{42}{52}\right)^2$</td>
</tr>
<tr>
<td>$(1, 1, 0)$</td>
<td>$\omega_1^1 \cdot \omega_2^1 \cdot \omega_3^0 = \left(\frac{27}{52}\right)^1 \left(\frac{25}{52}\right)^1 \left(\frac{19}{52}\right)^0 \left(\frac{42}{52}\right)^0$</td>
</tr>
<tr>
<td>$(1, 0, 1)$</td>
<td>$\omega_1^1 \cdot \omega_2^0 \cdot \omega_3^1 = \left(\frac{27}{52}\right)^1 \left(\frac{25}{52}\right)^0 \left(\frac{19}{52}\right)^1 \left(\frac{42}{52}\right)^1$</td>
</tr>
<tr>
<td>$(0, 1, 1)$</td>
<td>$\omega_1^0 \cdot \omega_2^1 \cdot \omega_3^1 = \left(\frac{27}{52}\right)^0 \left(\frac{25}{52}\right)^1 \left(\frac{19}{52}\right)^1 \left(\frac{42}{52}\right)^1$</td>
</tr>
</tbody>
</table>

we're using:

$$\omega_j^i = \left(\frac{27}{52}\right)^i \left(\frac{25}{52}\right)^j \left(\frac{19}{52}\right)^{j-i}$$

$$\omega_j^i = \left(\frac{19}{52}\right)^i \left(\frac{33}{52}\right)^j \left(\frac{42}{52}\right)^{j-i}$$

$$\omega_j^i = \left(\frac{42}{52}\right)^i \left(\frac{10}{52}\right)^j \left(\frac{52}{52}\right)^{j-i}$$
Closed Jackson Networks: Example

Example: Roommate network, continued

\[
\sum_{j_1, j_2, j_3 \geq 0} \omega_{j_1}^1 \cdot \omega_{j_2}^2 \cdot \omega_{j_3}^3 = \omega_0^0 \cdot \omega_0^2 \cdot \omega_0^3 = \left( \frac{27}{52} \right) \left( \frac{25}{52} \right) \left( \frac{19}{52} \right) \left( \frac{42}{52} \right) \\
+ \omega_0^1 \omega_0^2 \omega_0^3 = \left( \frac{27}{52} \right) \left( \frac{19}{52} \right) \left( \frac{33}{52} \right) \left( \frac{42}{52} \right) \\
+ \omega_0^1 \omega_0^2 \omega_0^3 = \left( \frac{27}{52} \right) \left( \frac{19}{52} \right) \left( \frac{42}{52} \right) \left( \frac{10}{52} \right) \\
+ \omega_0^1 \omega_0^2 \omega_0^3 = \left( \frac{27}{52} \right) \left( \frac{19}{52} \right) \left( \frac{33}{52} \right) \left( \frac{42}{52} \right) \\
+ \omega_0^1 \omega_0^2 \omega_0^3 = \left( \frac{27}{52} \right) \left( \frac{19}{52} \right) \left( \frac{42}{52} \right) \left( \frac{10}{52} \right) \\
\]

= 0.1824

Closed Jackson Networks: Example

Example: Roommate network, continued

\[
P_{n_1,n_2,n_3} = \begin{cases} 
0 & \text{if } n_1 + n_2 + n_3 \neq 2 \\
\sum_{j_1, j_2, j_3 \geq 0} \omega_{j_1}^1 \cdot \omega_{j_2}^2 \cdot \omega_{j_3}^3 & \text{if } n_1 + n_2 + n_3 = 2
\end{cases}
\]

The probability that each of your roommates has a Gameboy and you don’t is

\[
P_{0,1,1} = \frac{\sum_{j_1, j_2, j_3 \geq 0} \omega_{j_1}^1 \cdot \omega_{j_2}^2 \cdot \omega_{j_3}^3}{0.1824} = 0.0187 = 0.1025
\]
Queuing Models with Nonexponential Service Distribution

The exponential distribution is a very convenient choice for modeling service and interarrival time distributions. As we have seen, it leads to very tractable results in characterizing system behavior. In some contexts, however, it is not a realistic choice of distribution. In some service systems, the service time might be known to have a different distribution. One important example of this situation is when the service times are completely predictable (i.e., deterministic). We will consider the following types of non-exponential queues.

- M/G/1 queue
- M/D/1 queue
- M/E_k/1 queue

Nonexponential Service: The M/G/1 Queue

- Poisson input process with rate $\lambda$
- Single server
- Arbitrary service time distribution (iid for all customers) with mean $1/\mu$ and variance $\sigma^2$
- Infinite capacity queue

For this class of queues, if $\rho=\lambda/\mu<1$, then $P_0 = 1 - \rho$ and $L_Q = \frac{\lambda^2 \sigma^2 + \rho^2}{2(1 - \rho)}$

**Pollaczek-Khintchine Equation for M/G/1**

For an M/G/1 queue,

$$L_Q = \frac{\lambda^2 \sigma^2 + \rho^2}{2(1 - \rho)} \quad \text{if } \rho=\lambda/\mu<1.$$

Additional quantities of interest ($W$, $W_Q$, $L$, $L_s$) can be derived from $L_Q$ and $P_0$.
The M/G/1 Queue: Example

**EXAMPLE: McD’s Drive Thru**

Suppose that customers arrive at the McDonald’s drive-thru at a rate of 30 per hour. The time to until any given car completes service is uniformly distributed on the interval [0, 3] minutes.

1. What is the average amount of time a car spends waiting to be served?
2. What is the average number of cars being served or waiting to be served?

The mean service time is 1.5 minute (obtained by taking the mean of a uniform [0, 3] random variable) so the average service rate $\mu$ is $2/3$ cars/minute. Thus $\rho = \lambda / \mu = 3/4$. The variance of the service time is $\sigma^2 = (3)^2 / 12 = 3/4$. The average waiting time in queue for each car, measured in minutes, is computed using the P-K equation and our useful relationship between $L_Q$ and $W_Q$:

$$W_Q = \frac{L_Q}{\lambda} = \frac{\lambda \sigma^2 + \rho^2}{2 \lambda (1 - \rho)} = \frac{(0.5)^2 (0.75) + (0.75)^2}{2(0.5)(0.25)} = 3 \text{ minutes}$$

The M/G/1 Queue: Example

**EXAMPLE: McD’s Drive Thru**

The average number of cars being served or waiting to be served is

$$L = L_Q + L_s$$

Since $L_s = (1 - P_0) = \rho$ for an M/G/1 queue, $L$ equals

$$L = L_Q + L_s = \frac{\lambda \sigma^2 + \rho^2}{2(1 - \rho)} + \rho$$

$$= \frac{(0.5)^2 (0.75) + (0.75)^2}{2(0.25)} + (0.75) = 2.25 \text{ cars}$$
Nonexponential Service:
The M/ D/ 1 Queue

- Poisson input process with rate \( \lambda \)
- single server
- service time is deterministic and equal to \( 1/\mu \)
- infinite capacity queue

This is a special case of the M / G / 1 queue in which \( \sigma^2=0 \). As in the general case, if \( \rho=\lambda/\mu<1 \), then \( \rho_s = 1 - \rho \). Moreover, in this case the P-K equation reduces to

\[
L_q = \frac{\rho^2}{2(1-\rho)}
\]

**POISSON-KHINTCHINE EQUATION for M/ D/ 1**

For an M / D/ 1 queue,

\[
L_q = \frac{\rho^2}{2(1-\rho)} \quad \text{where} \quad \rho = \frac{\lambda}{\mu}
\]

The M/ D/ 1 Queue: Example

**EXAMPLE: McD's Drive Thru, again**

Let's compare two different employees at the drive thru window. Ann's service time is uniformly distributed on the interval \([0,3]\) minutes. Jim is slower on average, but more consistent: he completes service for every car in exactly 1.55 minutes every time. Which server makes customers wait less time from entry in the queue until completing service?

We've already seen that Ann's service performance leads to:

\[
W_q = \frac{L_q}{\lambda} = \frac{L_q}{\lambda} = \frac{\lambda^2 \sigma^2 + \rho^2}{2 \lambda (1-\rho)} = \frac{(0.5)^2(0.75) + (0.75)^2}{2(0.5)(0.25)} = 3 \text{ minutes}
\]

So

\[
W = W_q + W_s = 3 + 1.5 = 4.5 \text{ minutes}
\]
The M/ D/ 1 Queue: Example

**EXAMPLE: McD’s Drive Thru, continued**

Since Jim’s service time is deterministic, the variance of his service time is zero. In his case, $\rho = \lambda/\mu = 0.775$. By contrast, Jim has people wait in queue on average

$$W_q = \frac{L_q}{\lambda} = \frac{L_q}{\mu} = \frac{\rho^2}{2\lambda(1-\rho)} = \frac{(0.775)^2}{2(0.5)(0.225)} = 2.6694 \text{ minutes}$$

$$W = W_q + W_s = 2.6694 + 1.55 = 4.2194 \text{ minutes}$$

Although Ann’s average service time is faster, Jim’s consistency leads to shorter waits.

---

**Nonexponential Service:**

**The M/ E_k/ 1 Queue**

Now we consider another special case of the M/ G/ 1 queue, namely the M/ E_k/ 1 queue. The notation $E_k$ stands for the Erlang distribution with shape parameter $k$. To motivate this type of queue, we first mention some important facts about the Erlang distribution.

**The Erlang Distribution**

A random variable $X$ having an Erlang distribution with parameters $(\mu, k)$ has probability density function

$$f(x) = \frac{(\mu k)^k x^{k-1} e^{-\mu x}}{(k-1)!} \text{ for } x \geq 0$$

and mean and variance

$$E(X) = \frac{1}{\mu} \quad \text{Var}(X) = \frac{1}{k\mu^2}$$

The Erlang $(\mu, k)$ random variable is also called a Gamma $(k, \mu)$ random variable.

Note: if $k = 1$, then $X$ is exponentially distributed with parameter $\mu$. 
The Erlang Distribution

**Useful Fact About the Erlang Distribution**

Let \( Y_1, Y_2, \ldots, Y_n \) be independent and identically distributed exponential random variables with parameter \( n\mu \). Then the random variable

\[
Y = Y_1 + Y_2 + \ldots + Y_n
\]

has an Erlang distribution with parameters \((\mu, n)\).

The Erlang distribution is of particular significance in queueing theory because of this useful fact. To see why, imagine a queueing system where the server (or servers) performs not just a single task that takes an exponential length of time but instead a sequence of \( n \) tasks, where each task takes an exponential length of time. When each task’s time has parameter \( n\mu \) (mean \( 1/n\mu \)) then the length of time to complete all tasks has an Erlang distribution with parameters \((\mu, n)\).

Nonexponential Service:
The M/E\(_k\)/1 Queue

- Poisson input process with rate \( \lambda \)
- Single server
- Erlang (\( \mu, k \)) service time distribution; mean \( 1/\mu \) and variance \( 1/k\mu^2 \)
- Infinite capacity queue

As for the general M/G/1 case, if \( \rho = \lambda/\mu < 1 \), then \( P_0 = 1 - \rho \).

Applying the P-K equation to this case yields

\[
L_0 = \frac{\lambda^2 / k\mu^2 + \rho^2}{2(1 - \rho)} = \frac{\rho^2((1/k) + 1)}{2(1 - \rho)}
\]

**POLLACZEK-KHINTCHINE EQUATION for M/E\(_k\)/1**

For an \( M/E/1 \) queue,

\[
L_0 = \frac{\rho^2((1/k) + 1)}{2(1 - \rho)} \quad \text{if } \rho = \lambda/\mu < 1.
\]
The M/\(E_k/1\) Queue: Example

**EXAMPLE: Homework Questions**

Students working on their EESOR 121 homework send email asking for help at a rate of 1 email per hour. When an email arrives I get right to work answering that person's email if I am not busy answering another student's email at the time. Each student's email asks one question for each of the 7 problems on the homework.) The time each question takes me to answer is exponentially distributed with mean 3 minutes, so I am capable of replying to one email every 21 minutes on average.

(1) What is the average time a student waits for a reply?

(2) If I would like to spend only 10% of my time answering homework questions, how many homework problems should I assign each week?

**EXAMPLE: Homework, continued**

It is given that the time to answer each question is exponentially distributed with parameter 3. Since there are \(n=7\) questions on each homework set, letting \(1/n\mu=3\) implies that the time it takes me to answer one entire email has an Erlang distribution with parameters \((\mu, n)=(1/21, 7)\).

In this example \(\rho=\lambda/\mu = 7/20\). The P-K equation tells us

\[
L_\varphi = \frac{\rho^2 ((1/n) + 1)}{2(1 - \rho)} = \frac{(7/20)^2 ((1/7) + 1)}{2(13/20)} = \frac{7}{65}
\]

We also know \(L_s = 1 - P_\varnothing = \rho = \frac{7}{20}\)

(1) The average time a student waits for a reply is

\[
W = \frac{L}{\lambda} = \frac{L_s + L_\varphi}{\lambda} = \frac{7}{20} + \frac{7}{65} = 60\left(\frac{7}{20} + \frac{7}{65}\right) = 60\left(\frac{119}{260}\right)
\]

\[= 27.46 \text{ minutes} = 0.4577 \text{ hours}\]
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**EXAMPLE: Homework, continued**

(2) Let $n$ be the number of questions on the homework. My rate of replying to emails is then $\mu = 1/3n$ per hour. If I would like to spend at most 10% of my time answering homework questions, then I require

$$1 - P_0 \leq 0.1$$

This holds only if

$$1 - P_0 = \rho = \frac{\lambda}{\mu} = \frac{3n}{60} = \frac{n}{20} \leq 0.1$$

or

$$n \leq 2$$

I should assign at most 2 problems per week.